



THE UNIVERSITY of EDINBURGH
informatics

Security in Open Multi-agent Systems

Shahriar Bijani

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Outline

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- **Open Multi-agent Systems**
- **Attack Classification**
- **Probing Attack**
- **Attack Detection**

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Introduction: Open MAS

- **Open Multi-agent System (MAS)**

An open system in which autonomous agents can join and leave freely.

- **In Our Analysis:**

Open peer to peer MAS that agents can invent protocols for different applications and share them.

Introduction: LCC

Role	<u>a(requester, A) ::</u>	}	Clause		
Message out	<u>ask(X) ⇒ a(informer, B) ←</u>				
Constraint	<u>query_from(X, B) then</u>				
Message in	<u>tell(X) ← a(informer, B) then</u>			}	Role definition
Recursion	<u>a(requester, A)</u>				
	 a(informer, B) :: ask(X) ← a(requester, A) then tell(X) ⇒ a(requester, A) ← know(X)				

LCC language syntax:

Outgoing message: ⇒

Incoming message: ←

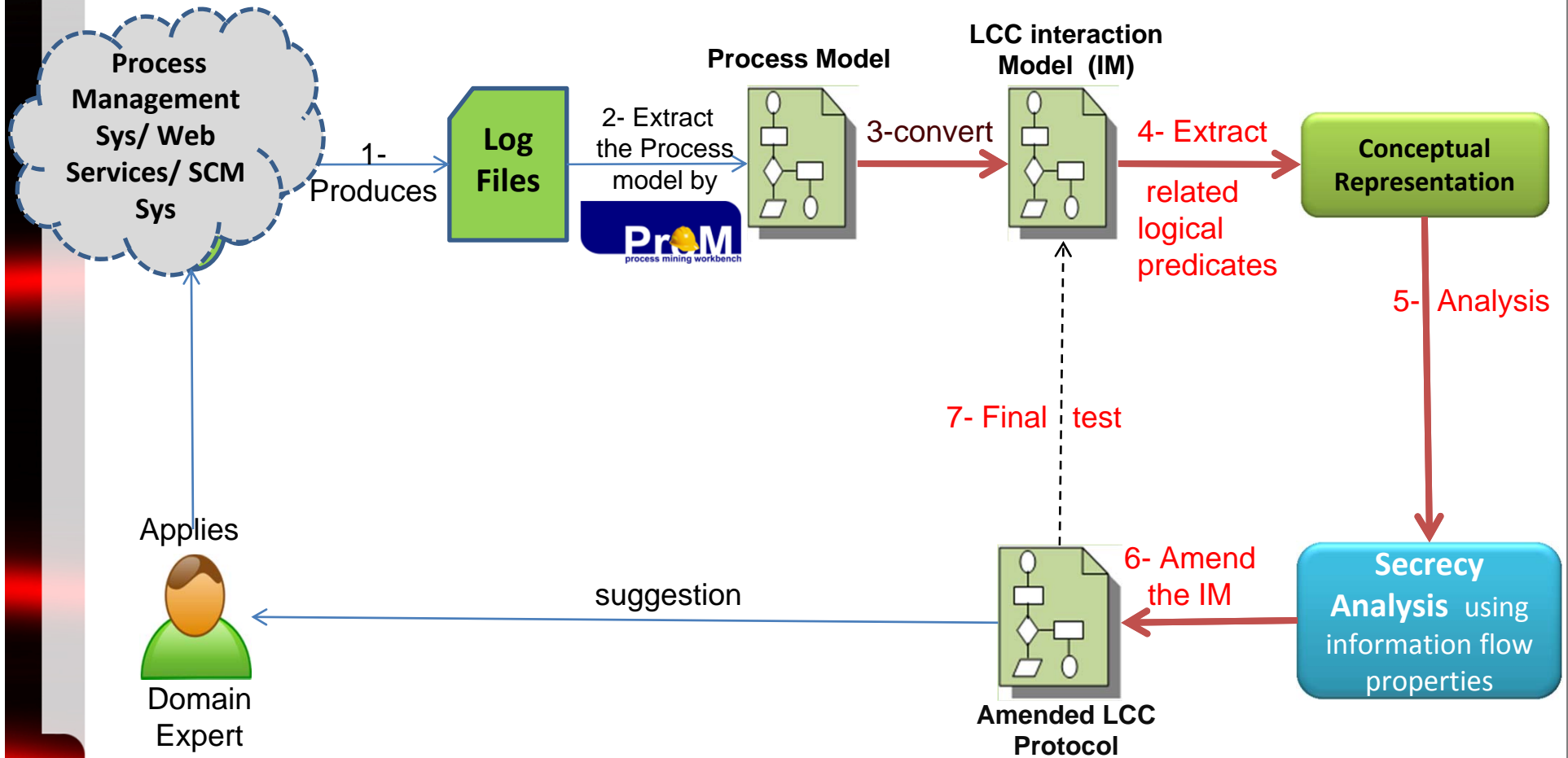
Conditional: ←

Sequence: **then**

Committed choice: **or**

- We might extend the scope of our security analysis without much difficulty to other domains such as web services.

Our Security Analysis Framework



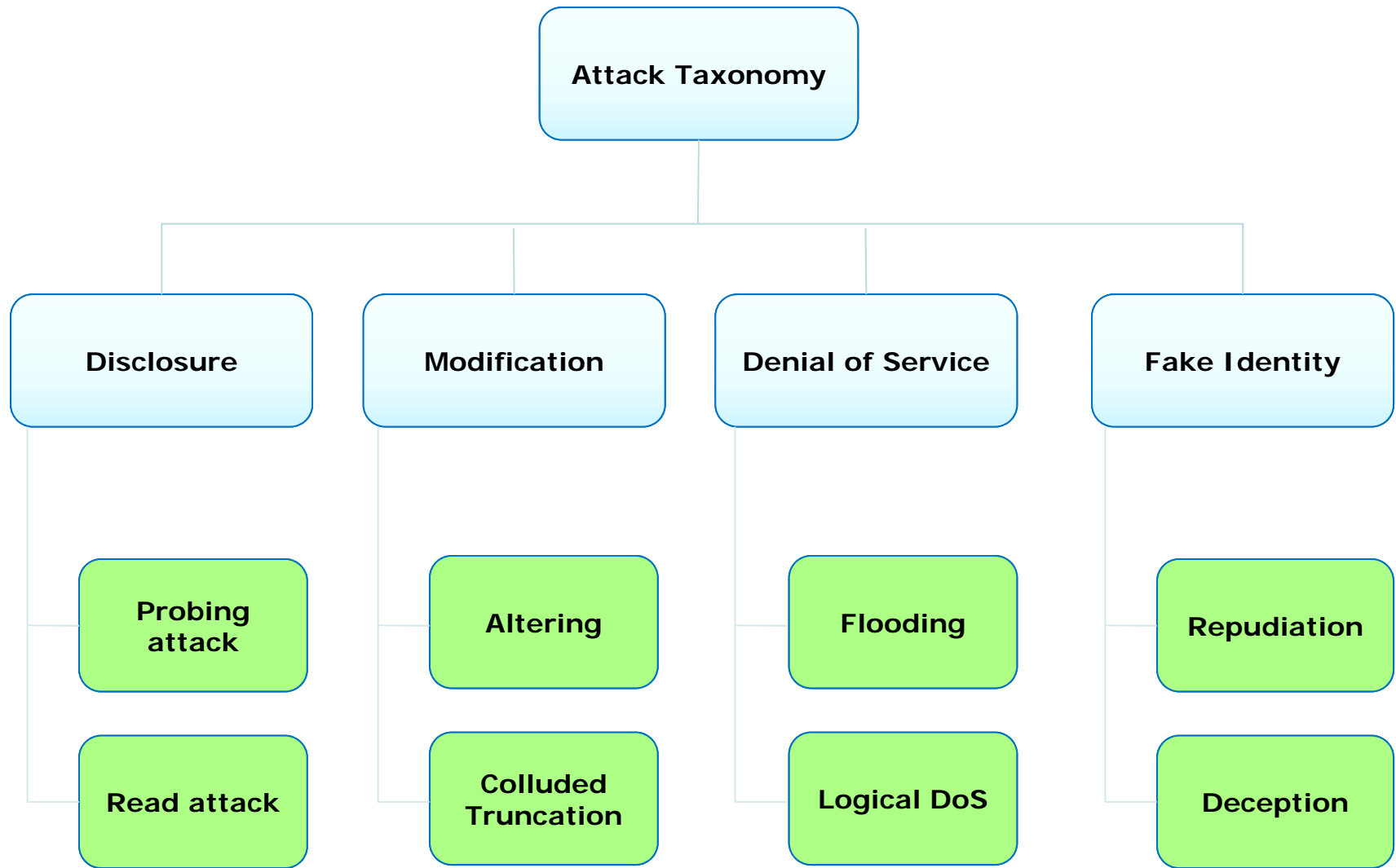
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Attack Classification



Attacks on Open MAS

Disclosure

- To intercept data, code, protocols, ...
- Tracing the interaction models

Modification

- To alter transferring information
- Malicious agents collude to modify

Fake Identity

- An attacker plays many pseudonymous roles
- A victim may enter into a fake interaction
- A method to deceive the *trust service*

Denial of Service

- To prevent agents to use provided service
- To waste others' resources
- To ruin the reputation

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Probing Attack

- an adversary could infer information not only from the interaction model itself, but also from the local knowledge of other agents.
- Four types of probing attacks:
 1. explicit query
 2. implicit query
 3. injection attack
 4. indirect query

Implicit Probing Attack

An Example of Proteomics Lab Interaction Model

```
a(researcher(LabList), R) ::  
  ( ask(X)=> a(omicslab, H) <-  
      LabList=[H|T] then  
    tell(Y)<= a(omicslab,H) then  
    null <- processResult(X,Y,H)  
    then a( researcher(T), R)  
  ) or  
  null <- LabList = []
```

```
a(omicslab, O) ::  
  ask(X)<= a(researcher,R) then  
  tell(Y)=>a(researcher,R)<-  
      Combine(X,Y)  
  then a(omicslab, O)
```

Injection Attack

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Injection Attack

- A Selling Interaction Model:

```
a(vendor, V)::  
  null <- ask(S) <= a(customer, C) then  
  null <- (not(want(C, S)) or payFor(C, S) ) then /*injection */  
  null <- (not(SupplyFrom(X)) or want(C, S) ) then /* injection */  
  ok => a(customer, C) <- agree(C, S)      then /* implicit query */  
  ...  
  then a(vendor, V)
```


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- **Attack Detection**

Injection Attack

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  ok => a(customer, C) <- agree(C, S)      then /* implicit query */  
  ...  
  then a(vendor, V)
```

- Conceptual representation:

```
IM1 = {  
  want(C, S) → payFor(C, S) ,  
  SupplyFrom(X) → want(C, S)  
}
```

```
q1 = {agree(C, S)}
```

Query result: **send(C)**

Attack Detection

- We formulate the problem as $A_0 \cup \text{IM}_i \vdash q$

Where:

A_0 = Agent's Local Knowledge

IM_i = The injection parts of an Interaction Model

q = query

- Example: $A_0 \cup \text{IM}_1 \vdash q_1$

$$A_0 \cup \text{IM}_2 \vdash q_1$$

$$A_0 \cup \text{IM}_3 \vdash q_1$$

Attack Detection

- **Detectability** (or non-opacity): an information flow property that shows the ability to infer a specific predicate from a set of rules
- We use an inference system for *detectability* [Becker 2010]

$$(PEEK) \frac{IM \vdash q \quad q \text{ is monotonic, } IM \text{ is ground}}{\Pi, IM \Vdash q}$$

$$(WEAK) \frac{\Pi, IM \Vdash q \quad q \Rightarrow q'}{\Pi, IM \Vdash q'}$$

$$(POKE1) \frac{A_0 \cup IM \vdash q \quad (IM, q) \in P}{\Pi, IM \Vdash q}$$

$$(POKE2) \frac{A_0 \cup IM \not\vdash q \quad (IM, q) \in P}{\Pi, IM \Vdash \neg q}$$

$$(MONO1) \frac{\Pi, IM \Vdash q \quad IM' \succcurlyeq IM \quad q \text{ is monotonic} \quad IM' \text{ is ground}}{\Pi, IM' \Vdash q}$$

$$(MONO2) \frac{\Pi, IM \Vdash q \quad IM' \preccurlyeq IM \quad \neg q \text{ is monotonic} \quad IM' \text{ is ground}}{\Pi, IM' \Vdash q}$$

$$(CONJ) \frac{\Pi, IM \Vdash q_1 \quad \Pi, IM \Vdash q_2}{\Pi, IM \Vdash q_1 \wedge q_2}$$

$$(DIFF) \frac{\Pi, A_0 \cup IM \vdash q}{\Pi, A_0 \Vdash q \vee \bigvee_{a \in IM} \text{fired}(IM, \text{head}(a))}$$

Countermeasures

- Two reasons that security problems lead to probing attacks:
 - (1) no distinguishing notion of private and public data in LCC
 - (2) no mechanism for information leakage control in an interaction.

Countermeasures

- Prevention
 - change the syntax of LCC
- Detection
 - An intrusion detection interaction model to interpret other interaction models

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