

## WEEK 6

We may draw out a tree structure with a formula at each node, in order to test a formula (or set of formulas) for consistency or inconsistency.

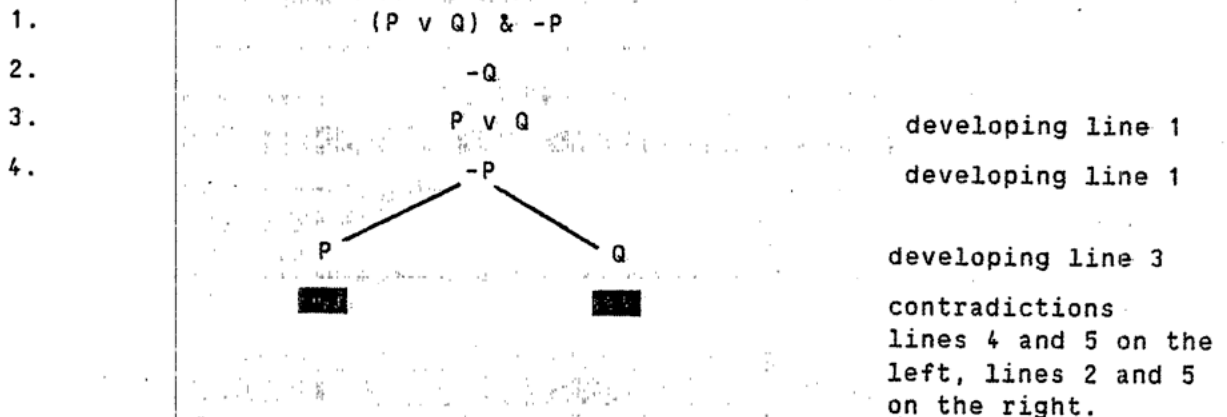
Begin by writing down the formula(s) to be tested, in a vertical column if there are two or more.

Each node (formula) in the tree other than an atom or a negated atom must be developed in every branch leading down from it.

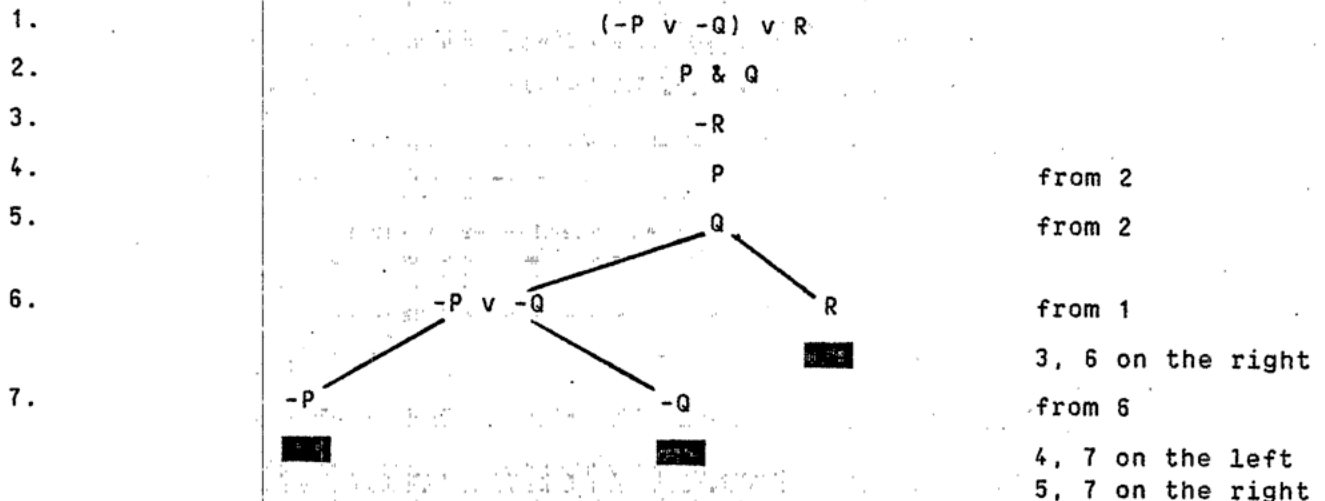
A conjunction  $A \& B$  is developed by writing the conjuncts one below the other. This represents the point that for the conjunction to be true both the conjuncts must be true as well. A disjunction  $A \vee B$  is developed by splitting the branch into two, one continued with  $A$  and the other with  $B$ . So the disjunction survives iff one or other disjunct survives.

A branch of the tree dies if it contains a contradiction: a pair of nodes of the form  $A$  and  $\neg A$ . The tree dies iff every branch dies. In that case, there is no way the formula(s) being tested could (all) be true.

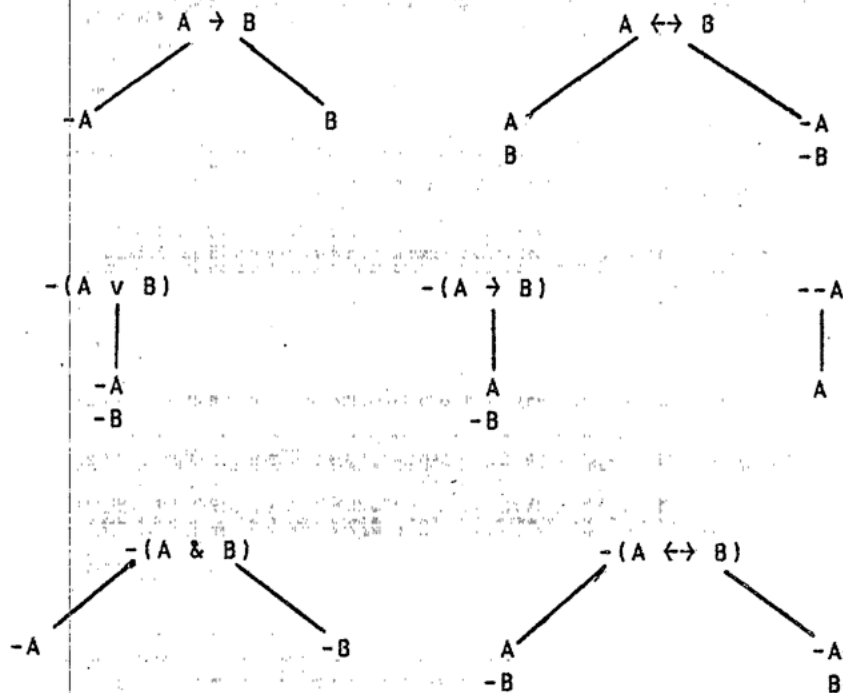
As an example, consider the two formulas  $(P \vee Q) \& \neg P, \neg Q$ . Could they both be true together? Let us see.



As another example, consider  $(\neg P \vee \neg Q) \vee R, P \& Q, \neg R$ .



# Rules for other connectives and negated cases



To test a sequent  $A_1, \dots, A_n : B$  for validity, see whether there is any way all the premisses could be true and the conclusion false. That is, test the set  $\{A_1, \dots, A_n, \neg B\}$  for consistency. Suppose all the premisses true and the conclusion,  $B$ , false. Develop the truth tree by the rules given above. If all branches die because of contradictions, the sequent is valid. For instance, let us test

$$P \rightarrow \neg(Q \& R), \neg(P \rightarrow \neg Q) : R \rightarrow \neg P$$

1.	$P \rightarrow \neg(Q \& R)$	premiss
2.	$\neg(P \rightarrow \neg Q)$	premiss
3.	$\neg(R \rightarrow \neg P)$	neg. conclusion
4.	$P$	from 2
5.	$\neg Q$	from 2
6.	$R$	from 3
7.	$\neg P$	from 3
8.	$\neg P$ (contradiction) $\neg(Q \& R)$	from 1 4 and 8
9.	$\neg Q$ (contradiction) $\neg R$ (contradiction)	from 8 5/6 and 9.

All the branches eventually contain contradictions, so the sequent is valid. To produce a proof from primitive rules in Lemmon's formulation of logic would be rather difficult, but the truth tree method handles it very easily.